

Dedekind-MacNeille completions of residuated lattices

Joint work with A. Ciabattoni and K. Terui

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- $(A, \wedge, \vee, 0, 1)$ is a bounded lattice,
- for all $a, b, c \in A$,

$$a \wedge b \leq c \Leftrightarrow b \leq a \rightarrow c \text{ (\wedge-residuation)}$$

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Fact: The DM-completion of \mathbf{A} is the unique (up to isomorphism) completion in which \mathbf{A} is both meet dense and join dense. Namely, every element a can be written as

$$a = \bigvee X = \bigwedge Y \quad \text{for some } X, Y \subseteq A.$$

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Proposition: Gödel algebras are closed under completions.

Fact: $A \models (x \rightarrow y) \vee (y \rightarrow x) = 1$ iff $A \models x \leq y$ or $y \leq x$ (*lin*), for $A \in \mathbf{HA}_{SI}$.

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This contradicts (*lin'*).

A *residuated lattice*, or *residuated lattice-ordered monoid*, is an algebra $\mathbf{L} = (L, \wedge, \vee, \cdot, \backslash, /, 1)$ such that

- (L, \wedge, \vee) is a lattice,
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Fact. The last condition is equivalent to either one of:

- Multiplication **distributes over existing \vee 's** and, for all $a, c \in L$, $\bigvee\{b : ab \leq c\}$ ($=: a \backslash c$) and $\bigvee\{b : ba \leq c\}$ ($=: c / a$) exist.

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Therefore, the class **RL** of residuated lattices is an **equational class/variety**. We write $x \rightarrow y$ for $x \backslash y$ and y / x , when they are equal.

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$$\begin{array}{ll} b \leq a \backslash (ab \vee c) & a \leq (c \vee ab) / b \\ a(a \backslash c \wedge b) \leq c & (a \wedge c / b)b \leq c \end{array}$$

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We also add in the language a constant 0 , for which we stipulate nothing. It allows the definition of negation(s) $\neg x := x \rightarrow 0$.

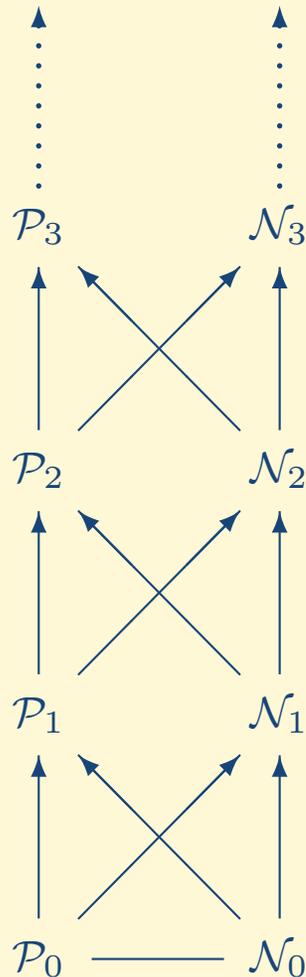
- **Lattice-ordered groups.** For $x \backslash y = x^{-1}y$, $y / x = yx^{-1}$.
- **(Reducts of) relation algebras.** For $x \cdot y = x; y$,
 $x \backslash y = (x^\cup; y^c)^c$, $y / x = (y^c; x^\cup)^c$, $1 = id$ and $0 = id^c$.
- The **powerset** $(\mathcal{P}(M), \cap, \cup, \cdot, \backslash, /, \{e\})$ of a monoid $\mathbf{M} = (M, \cdot, e)$, where $X \cdot Y = \{x \cdot y \mid x \in X, y \in Y\}$,
 $X / Y = \{z \in M \mid \{z\} \cdot Y \subseteq X\}$,
 $Y \backslash X = \{z \in M \mid Y \cdot \{z\} \subseteq X\}$.
- **Ideals of a ring (with 1)**, where $IJ = \{\sum_{fin} ij \mid i \in I, j \in J\}$
 $I / J = \{k \mid kJ \subseteq I\}$, $J \backslash I = \{k \mid Jk \subseteq I\}$, $1 = R$.
- **Quantales** are (essentially) complete residuated lattices.
- **Boolean algebras.** $x / y = y \backslash x = y \rightarrow x = y^c \vee x$ and
 $x \cdot y = x \wedge y$.
- **MV-algebras.** For $x \cdot y = x \odot y$ and $x \backslash y = y / x = \neg(\neg x \odot y)$.
- Models of **relevance** and of **linear logic**.

For $\{\vee, \cdot, 1\}$

- $x \cdot 1 = x = 1 \cdot x$
- $x(y \vee z) = xy \vee xz$ and $(y \vee z)x = yx \vee zx$

For $\{\wedge, \backslash, /\}$ (and $\{\vee, \cdot, 1\}$ in the denominator)

- $x \backslash (y/z) = (x \backslash y)/z$
- $x \backslash (y \wedge z) = (x \backslash y) \wedge (x \backslash z)$ and $(y \wedge z)/x = (y/x) \wedge (z/x)$
- $(y \vee z) \backslash x = (y \backslash x) \wedge (z \backslash x)$ and $x/(y \vee z) = (x/y) \wedge (x/z)$
- $(yz) \backslash x = z \backslash (y \backslash x)$ and $x/(zy) = (x/y)/z$
- $1 \backslash x = x = x/1$



- Polarity $\{\vee, \cdot, 1\}, \{\wedge, \backslash, /\}$
- The sets $\mathcal{P}_n, \mathcal{N}_n$ of terms are defined by:
 - (0) $\mathcal{P}_0 = \mathcal{N}_0 =$ the set of variables
 - (P1) $\mathcal{N}_n \cup \{1\} \subseteq \mathcal{P}_{n+1}$
 - (P2) $\alpha, \beta \in \mathcal{P}_{n+1} \Rightarrow \alpha \vee \beta, \alpha \cdot \beta \in \mathcal{P}_{n+1}$
 - (N1) $\mathcal{P}_n \cup \{0\} \subseteq \mathcal{N}_{n+1}$
 - (N2) $\alpha, \beta \in \mathcal{N}_{n+1} \Rightarrow \alpha \wedge \beta \in \mathcal{N}_{n+1}$
 - (N3) $\alpha \in \mathcal{P}_{n+1}, \beta \in \mathcal{N}_{n+1} \Rightarrow \alpha \backslash \beta, \beta / \alpha \in \mathcal{N}_{n+1}$
- $\mathcal{P}_{n+1} = \langle \mathcal{N}_n \rangle_{\vee, \Pi}$; $\mathcal{N}_{n+1} = \langle \mathcal{P}_n \rangle_{\wedge, \mathcal{P}_{n+1} \backslash, / \mathcal{P}_{n+1}}$
- $\mathcal{P}_n \subseteq \mathcal{P}_{n+1}, \mathcal{N}_n \subseteq \mathcal{N}_{n+1}, \bigcup \mathcal{P}_n = \bigcup \mathcal{N}_n = Fm$
- \mathcal{P}_1 -reduced: $\bigvee \prod p_i$
- \mathcal{N}_1 -reduced: $\bigwedge (p_1 p_2 \cdots p_n \backslash r / q_1 q_2 \cdots q_m)$

\mathcal{N}_2 -normal formulas are of the form $\alpha_1 \cdots \alpha_n \rightarrow \beta$ where

- $\beta = 0$ or $\beta_1 \vee \cdots \vee \beta_k$ with each β_i a product of variables
- each α_i is of the form $\bigwedge_{1 \leq j \leq m_i} \gamma_i^j \rightarrow \beta_i^j$, where $\beta_i^j = 0$ or is a variable, and γ_i^j is a product of variables.

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For any set E of \mathcal{N}_2 -equations, the following are equivalent:

- The variety $Mod(E)$ is closed under completions.
- The variety $Mod(E)$ is closed under DM-completions.
- E is equivalent to a set of acyclic equations.
- E is equivalent to a set of analytic equations.

If E implies integrality $x \leq 1$, all the above hold.

- Heyting algebras
- Gödel algebras
- Why?
- Residuated lattices
- Examples
- Properties
- Term hierarchy
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A *structural clause* is a universal first-order formula of the form:

$$t_1 \leq u_1 \text{ and } \dots \text{ and } t_m \leq u_m \implies t_{m+1} \leq u_{m+1} \text{ or } \dots \text{ or } t_n \leq u_n$$

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Let $L = \text{var}\{t_{m+1}, \dots, t_n\}$ and $R = \text{var}\{u_{m+1}, \dots, u_n\}$. The clause is called *analytic* if it satisfies:

- L and R are disjoint.
- Each variable occurs only once in $t_{m+1}, u_{m+1}, \dots, t_n, u_n$.
- $\text{var}\{t_1, \dots, t_m\} \subseteq L$ and $\text{var}\{u_1, \dots, u_m\} \subseteq R$.

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In the absence of commutativity and integrality we need to consider:

1. *iterated conjugates*:

A *conjugate* of a term t is either $\lambda_u(t) = (u \setminus tu) \wedge 1$ or $\rho_u(t) = (ut / u) \wedge 1$ for some term u . We have:

$$\lambda_u(t) \leq 1, \quad \rho_u(t) \leq 1, \quad u\lambda_u(t) \leq tu, \quad \rho_u(t)u \leq ut.$$

Iterated conjugates are compositions of conjugates.

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2. *acyclic clauses*:

A clause

$$t_1 \leq u_1 \text{ and } \dots \text{ and } t_m \leq u_m \implies t_{m+1} \leq u_{m+1} \text{ or } \dots \text{ or } t_n \leq u_n$$

is called *acyclic* if there are no directed cycles in the directed graph (G, E) , where $G = \text{var}\{t_1, u_1, \dots, t_m, u_m\}$, and $(x, y) \in E$ iff $\text{lxr} \leq y$ is a premise.

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Theorem. Every acyclic structural clause is equivalent to an analytic one.

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